

# COUNTING WORD-TREES

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**Introduction.** In his recent study of free inverse semigroups, Munn [2] introduced and used extensively the concept of a *word-tree*. In this note the number of such trees is found.

**DEFINITION.** A word-tree on an alphabet  $A$  is a finite tree, with at least two points, that satisfies the following conditions.

(WT1) Each line is oriented and is labelled by an element of  $A$ .

(WT2)  $T$  has no subgraph of the form  $o \xrightarrow{a} o \xleftarrow{a} o$  or  $o \xleftarrow{a} o \xrightarrow{a} o$  ( $a \in A$ ).

Lines will be described by an ordered pair of adjacent points and are said to be *similar* if their orientations and labellings are the same.

A line  $\alpha\beta$ , where  $\beta$  is an endpoint of a word-tree  $T$ , is called an *endline* of  $T$ .

Isomorphism is defined in the obvious way (preserving orientation and labelling of lines).

A fundamental fact proved in Munn's paper [2, Theorem 2.2] is that a word-tree has no nontrivial automorphism.

*Notations.*  $d(\alpha)$  = degree of point  $\alpha$ .

$k$  = cardinal of  $A$ , taken to be finite.

$n(p)$  = number of word-trees on  $A$  with  $p$  points.

$n(p, e)$  = number of word-trees on  $A$  with  $p$  points and  $e$  endpoints.

$$(n)_j = \begin{cases} n(n-1)\dots(n-j+1) & (j = 1, 2, \dots), \\ 1 & (j = 0). \end{cases}$$

Note that  $n(p, e) = 0$  if  $e < 2$  or  $p < e$  and that, for  $p > 2$ ,  $n(p, p) = 0$ .

**Construction.** Given a word-tree  $T$  with  $p-1$  points and  $e$  endpoints, we can construct a larger one by adding a new line at any point. The number of dissimilar lines that may be added at  $\alpha$  is  $2k-d(\alpha)$ , since each old line terminating in  $\alpha$  imposes one restriction on the new line and, by (WT2), these restrictions are different. The new word-tree has  $e$  endpoints if the new line was added at an old endpoint and  $e+1$  endpoints otherwise.

Let  $T_1$  and  $T_2$  be the word-trees obtained by adding endlines  $\alpha_1\beta_1$  and  $\alpha_2\beta_2$  respectively to  $T$ , where  $\alpha_1$  and  $\alpha_2$  are distinct points of  $T$ . Then, since  $T$  has no nontrivial automorphisms, there can be no isomorphism  $\phi: T_1 \rightarrow T_2$  such that  $\beta_1\phi = \beta_2$ . Clearly the same result holds if  $\alpha_1 = \alpha_2$  but the endlines  $\alpha_1\beta_1$  and  $\alpha_2\beta_2$  are dissimilar. It follows that there is a 1-1 correspondence between the possible constructions on word-trees with  $p-1$  points and the word-trees on  $p$  points with a distinguished endpoint.

**Calculation.** From the tree  $T$  the construction gives  $(2k-1)e$  word-trees with  $e$  endpoints and  $\sum_{\alpha} (2k-d(\alpha))$  with  $e+1$  endpoints, where the sum is over all points  $\alpha$  of  $T$  that are not

endpoints of  $T$ . But  $\sum_{\alpha} d(\alpha) + e = 2$  (number of lines in  $T$ ) =  $2(p-2)$  (see, e.g., [1, Theorem 4.11]), so that the number of word-trees with  $e+1$  endpoints is

$$2k(p-1-e) - 2(p-2) + e.$$

Among all possible constructions, a particular tree with  $p$  points and  $e$  points will occur  $e$  times. Thus

$$\begin{aligned} n(p, e) &= \frac{1}{e} \{ (2k-1)e n(p-1, e) + (2k(p-1-e+1) - 2(p-2) + e - 1) n(p-1, e-1) \} \\ &= (2k-1)n(p-1, e) + \left\{ \frac{2kp-2p+3}{e} - (2k-1) \right\} n(p-1, e-1), \quad \text{for } p \geq 3. \end{aligned} \tag{1}$$

This recurrence formula is used  $p-3$  times to express  $n(p)$ , where  $p \geq 3$ , in terms of  $n(3, 2)$ , which is  $k(2k-1)$ . In fact, for  $0 \leq j \leq p-3$ ,

$$n(p) = (2kp - 2p + j + 2)_j \sum_{e=2}^{p-j-1} n(p-j, e) / (e+j)_j. \tag{2}$$

This is proved by induction on  $j$ . Application of (1) in (2) gives

$$\begin{aligned} n(p) &= (2kp - 2p + j + 2)_j \sum_{e=2}^{p-j-1} \left[ (2k-1)n(p-j-1, e) + \right. \\ &\quad \left. + \left\{ \frac{2k(p-j) - 2p + 2j + 3}{e} - (2k-1) \right\} n(p-j-1, e-1) \right] / (e+j)_j. \end{aligned}$$

In the sum, the coefficient of  $n(p-j-1, e)$ , when  $2 \leq e \leq p-j-2$ , is

$$\frac{2k-1}{(e+j)_j} + \frac{2k(p-j) - 2p + 2j + 3}{(e+1)(e+1+j)_j} - \frac{2k-1}{(e+1+j)_j} = (2kp - 2p + j + 3) / (e+j+1)_{j+1}.$$

The other substitute terms are 0. This verifies (2).

In particular, for  $j = p-3$ , we obtain

$$\begin{aligned} n(p) &= (2kp - 2p + p - 1)_{p-3} \sum_{e=2}^2 n(3, e) / (e+p-3)_{p-3} \\ &= (2kp - p - 1)_{p-3} k(2k-1) / \frac{(p-1)!}{2} \\ &= \frac{2k(2k-1)}{(p-1)(p-2)} \binom{2kp-p-1}{p-3}. \end{aligned}$$

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REFERENCES

1. F. Harary, *Graph theory* (Reading, Mass., 1969).
2. W. D. Munn, Free inverse semigroups, *Proc. London Math. Soc.*; to appear.

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