

A COUNTEREXAMPLE TO A CONJECTURE OF D. B. FUKS

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Introduction. In [3] D. B. Fuks defined a duality of functors in the category \mathcal{H} of weak homotopy types. In general this duality is more difficult to work with than the duality of functors of the category \mathcal{T} of pointed Kelley spaces [2]. It happens however that all so-called strong functors [2] F of \mathcal{T} induce functors \bar{F} of \mathcal{H} , and if we denote the duality operators of \mathcal{H} and \mathcal{T} by \mathcal{D} and D respectively, then there are many cases where $(\overline{DF}) = \mathcal{D}(\bar{F})$.

This has lead Fuks to make the following conjecture: a functor F of the category \mathcal{H} is reflexive (i.e. $F \simeq \mathcal{D}^2 F$) if and only if there exists a functor G of \mathcal{T} such that $F = \bar{G}$ and $\mathcal{D}F = \overline{DG}$.

Not only would this conjecture enable us to compute $\mathcal{D}F$ in the most interesting cases, but it would imply the following strong corollary.

COROLLARY. *Let G_1 and G_2 be reflexive functors of the category \mathcal{T} , and let $f: G_1 \rightarrow G_2$ be a natural transformation such that for any C.W. complex A , $f_A: G_1(A) \rightarrow G_2(A)$ is a weak homotopy equivalence. Then for any C.W. complex A ,*

$$(Df)_A: DG_2(A) \rightarrow DG_1(A)$$

is also a weak homotopy equivalence.

Unfortunately, we will provide a counterexample to this corollary, which will prove that Fuks' conjecture is false.

1. The Counterexample. Consider the functor QX =space of paths in X which start or end at the base point.

In other words, QX is the pull-back of the diagram

$$\begin{array}{ccc} & (I', X) & \\ & \downarrow & \\ X \vee X & \rightarrow & X \times X \end{array}$$

where I' = disjoint union of the unit interval $[0,1]$ with a point $*$ serving as the base point. The horizontal map is the inclusion, and the vertical map is defined as $\lambda \rightsquigarrow (\lambda(0), \lambda(1))$.

In [4, p. 210] there is defined a natural transformation

$$m: \Sigma\Omega \rightarrow Q$$

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with the property that for any 1-connected C.W. complex X , $m_X: \Sigma\Omega X \rightarrow QX$ is a homotopy equivalence.

This implies in particular that

$$(m * \Sigma^2)_X: \Sigma\Omega\Sigma^2 X \rightarrow Q\Sigma^2 X$$

is a homotopy equivalence for any C.W. complex X .

Now if we assume that Q is a reflexive functor (which will be proved later), Fuks' conjecture would imply that

$$((Dm) * \Omega^2)_X: DQ \circ \Omega^2 X \rightarrow \Omega\Sigma\Omega^2 X$$

is a homotopy equivalence for all C.W. complexes X .

In particular, if we take $X = K(Z, 3)$, we have $\Omega^2 X \simeq S^1 = K(Z, 1)$, so that $DQ(S^1) \simeq \Omega\Sigma(S^1)$. But from the explicit computation of DQ , it will be clear that $DQ(S^1)$ is a finite complex, while the homology of $\Omega\Sigma(S^1)$ is infinite, by a theorem of Bott and Samelson [1]. Hence $DQ(S^1)$ and $\Omega\Sigma(S^1)$ cannot have the same homotopy type.

We will now introduce a new functor which will be proved to be the dual of Q .

2. **The Functor T .** QX was defined as the pull-back of the natural transformations

$$X \vee X \rightarrow X \times X$$

and

$$(I', X) \rightarrow X \times X.$$

Now if we take the dual of this situation, we get two transformations:

$$X \vee X \rightarrow X \times X$$

and

$$X \vee X \rightarrow X \wedge I'$$

where the first map is the same as the first one above, and the second is defined as

$$(x, *) \rightsquigarrow (x, 0)$$

$$(*, x) \rightsquigarrow (x, 1)$$

We define then TX to be the push-out of the diagram

$$\begin{array}{c} X \vee X \rightarrow X \times X \\ \downarrow \\ X \wedge I' \end{array}$$

Since the duality operator D transforms direct limits into inverse limits, it is clear that $DT \simeq Q$. Because of this, there is a natural transformation $\sigma: T \rightarrow DQ$ corresponding to the canonical transformation $T \rightarrow D^2T$. Explicitly, σ is defined as follows:

From the definition of Q as a pull-back, we have two natural transformations

$$\omega_X: QX \rightarrow X \vee X \quad (\omega_X(\lambda) = (\lambda(0), \lambda(1)))$$

and

$$\tau_X: QX \rightarrow (I', X) \quad (\tau_X(\lambda) = \lambda).$$

If we denote the wedge functor by W (i.e. $W(X) = X \vee X$), then there is a natural equivalence $\alpha_X: X \times X \rightarrow D(W)(X)$ (see [2]) given by the formula

$$\begin{aligned} \alpha_X(x_0, x_1)_Y(y, *) &= (x_0, y) \in X \wedge Y \\ \alpha_X(x_0, x_1)_Y(*, y) &= (x_1, y) \in X \wedge Y. \end{aligned}$$

Thus we have a natural map $X \times X \rightarrow DQX$ obtained as the composition

$$X \times X \xrightarrow{\alpha_X} D(W)(X) \xrightarrow{(D\omega)_X} DQ(X).$$

Explicitly,

$$\begin{aligned} [(D\omega)_X \circ \alpha_X(x_0, x_1)]_Y(\lambda) &= \alpha_X(x_0, x_1)_Y \omega_Y(\lambda) \\ &= \alpha_X(x_0, x_1)(\lambda(0), \lambda(1)) \\ &= (x_0, \lambda(0)) \quad \text{if } \lambda(1) = * \\ &= (x_1, \lambda(1)) \quad \text{if } \lambda(0) = *. \end{aligned}$$

On the other hand, by taking the dual of τ , we obtain a map $(D\tau)_X: X \wedge I' \rightarrow DQ(X)$ given by the formula $(D\tau)_X(x, t)_Y(\lambda) = (x, \lambda(t)) \in X \wedge Y$, where $\lambda \in QY$. It is easy to see that the two maps

$$X \times X \rightarrow DQX$$

and

$$X \wedge I' \rightarrow DQX$$

agree on $X \vee X$, so that they combine to give us a map

$$\sigma_X: TX \rightarrow DQX.$$

3. The Functor \bar{T} and the Map $\eta: DQ \rightarrow \bar{T}$. Instead of finding directly an inverse for σ , we will introduce an auxiliary functor \bar{T} , define $\eta: DQ \rightarrow \bar{T}$ and show first that $\eta \circ \sigma$ is a natural equivalence of functors.

As usual, I will stand for the unit interval $[0, 1]$ with 0 as the base-point; we will let J denote the unit interval $[0, 1]$ with 1 as the base point.

$\rho: I \rightarrow S^1$ and $\pi: J \rightarrow S^1$ will denote the usual identification maps. $\bar{T}X$ is then defined as the pull-back of the diagram

$$\begin{array}{ccc} & X \wedge I & \\ & \downarrow X \wedge \rho & \\ X \wedge J & \xrightarrow{X \wedge \pi} & \Sigma X \end{array}$$

Define $\mu_x: DQX \rightarrow X \wedge J$ as $\mu(S) = T_J(\text{id}_J)$, where $S: Q \rightarrow \Sigma_x$ is an element of DQX , and id_J denotes the identity map of J , and

$$\nu_x: DQX \rightarrow X \wedge I$$

as

$$\nu_x(S) = S_I(\text{id}_I).$$

These maps are continuous, since they are the composition of the following continuous maps:

$$(Q, \Sigma_x) \xrightarrow[\text{at } J]{\text{evaluation}} (QJ, X \wedge J) \xrightarrow[\text{at } \text{id}_J]{\text{evaluation}} X \wedge J$$

and

$$(Q, \Sigma_x) \xrightarrow[\text{at } I]{\text{evaluation}} (QI, X \wedge I) \xrightarrow[\text{at } \text{id}_I]{\text{evaluation}} X \wedge I.$$

Moreover, because of the naturality of $S: Q \rightarrow \Sigma_x$, we have

$$(X \wedge \rho) \circ \nu_x(S) = (X \wedge \pi) \circ \mu_x(S) = S_{S^{-1}(\rho)}.$$

Thus the maps μ_x and ν_x induce a unique $\eta_x: DQX \rightarrow \bar{TX}$.

We will now prove that $\eta \circ \sigma$ is a natural equivalence.

4. The Natural Equivalence $\eta \circ \sigma: T \rightarrow \bar{T}$. From the definition of T as a push-out, we have that

$$TX = X \wedge I' \bigcup_{X \vee X} X \times X.$$

Let $(x, t) \in X \wedge I'$. We have

$$\mu_x \circ \sigma_x(x, t) = \sigma_x(x, t)_J(\text{id}_J) = (x, t) \in X \wedge J$$

$$\nu_x \circ \sigma_x(x, t) = \sigma_x(x, t)_I(\text{id}_I) = (x, t) \in X \wedge I.$$

Hence

$$\eta_x \circ \sigma_x(x, t) = ((x, t), (x, t)) \in X \wedge J \times X \wedge I.$$

On the other hand,

$$\mu_x \circ \sigma_x(x_0, x_1) = \sigma_x(x_0, x_1)(\text{id}_J) = (x_0, 0) \in X \wedge J$$

$$\nu_x \circ \sigma_x(x_0, x_1) = \sigma_x(x_0, x_1)_I(\text{id}_I) = (x_1, 1) \in X \wedge I.$$

Thus $(\eta \circ \sigma)_x: TX \rightarrow \bar{TX}$ is the map defined as

$$\left. \begin{aligned} (\eta\sigma)_x(x, t) &= ((x, t), (x, t)) \\ (\eta\sigma)_x(x_0, x_1) &= ((x_0, 0), (x_1, 1)) \end{aligned} \right\} \in X \wedge J \times X \wedge I.$$

It is clear that $(\eta\sigma)_x$ is both an injection and a surjection, so that it remains to show that its inverse is continuous. To that end, we define two subspaces A and B of \bar{TX} with the properties that

- (1) A and B are closed in \bar{TX}
- (2) $\bar{TX} = A \cup B$
- (3) $(\eta\sigma)^{-1}$ is continuous on both A and B .

$$A = \{((x, t), (x, t)) \in X \wedge J \times X \wedge I\}$$

$$B = \{((x_0, 0), (x_1, 1)) \in X \wedge J \times X \wedge I\}$$

Note first that under the homeomorphism

$$X \wedge J \times X \wedge I \rightarrow X \wedge I \times X \wedge I$$

sending $((x, t), (y, s))$ to $((x, 1-t), (y, s))$, A and B are mapped respectively onto

$$A^1 = \{((x, 1-t), (x, t)) \in X \wedge I \times X \wedge I\}$$

and

$$B^1 = \{(x, 1) \times (x_1, 1) \in X \wedge I \times X \wedge I\}$$

Let $q: X \times I \rightarrow X \wedge I$ be the identification map, and consider

$$q \times q: (X \times I) \times (X \times I) \rightarrow (X \wedge I) \times (X \wedge I).$$

The inverse image of A^1 under $q \times q$ is the union of the four subspaces

$$\{((x, 1-t), (x, t)) \in (X \times I) \times (X \times I)\}$$

$$\{((x, 0), (y, 1)) \in (X \times I) \times (X \times I)\}$$

$$\{(x, 1), (y, 0) \in (X \times I) \times (X \times I)\}$$

$$(q \times q)^{-1}(*).$$

These are all closed subsets of $(X \times I) \times (X \times I)$. If we can show that $q \times q$ is an identification map, then we will have proved that A^1 is closed.

But in the category of nonpointed Kelley spaces, the product with a fixed space is left adjoint to a hom functor and hence commutes with direct limits. Thus we have two push-out diagrams:

$$\begin{array}{ccc} (X \times \{0\}) \times (X \times I) & \longrightarrow & (X \times I) \times (X \times I) \\ \downarrow & & \downarrow q' \\ * \times (X \times I) & \longrightarrow & (X \wedge I) \times (X \times I) \end{array}$$

and

$$\begin{array}{ccc} (X \wedge I) \times (X \times \{0\}) & \longrightarrow & (X \wedge I) \times (X \times I) \\ \downarrow & & \downarrow q'' \\ (X \wedge I) \times * & \longrightarrow & (X \wedge I) \times (X \wedge I) \end{array}$$

Since q' and q'' are both identification maps, so is their composition $q'' \circ q' = q \times q$.

Thus A^1 is closed. As for B^1 , it is the product of two closed subsets of $X \wedge I$.

It remains only to show that $(\eta \circ \sigma)^{-1}$ is continuous on both A and B .

For A , we have the following commutative diagram

$$\begin{array}{ccc} X \wedge I' & \xrightarrow{(\eta\sigma)} & (X \wedge J) \times (X \wedge I) \\ \uparrow p & & \uparrow r \times q \\ X \times I & \xrightarrow{\nabla} & (X \times I) \times (X \times I) \end{array}$$

where $p, q,$ and r are identification maps and ∇ is the map

$$\nabla(x, t) = ((x, 1-t), (x, t)).$$

Since ∇ is a homeomorphism into, $(r \times q) \circ \nabla$ is an identification map onto its image, which is A . Hence $\eta\sigma$ is a homeomorphism of $X \wedge I'$ onto A , because p is also an identification map.

As for B , it is clear that $\eta\sigma: X \times X \rightarrow B$ is a homeomorphism, because B has the topology of a product.

This insures the continuity of $(\eta\sigma)^{-1}$ over the whole of \bar{TX} .

5. σ And η are Natural Equivalence. We have just shown that $\eta \circ \sigma$ is an isomorphism of functors. But η is a monomorphism. Indeed, let $R, S: Q \rightarrow \Sigma_X$ be two elements of DQX . We have

$$\eta_x(R) = (R_I(\text{id}_I), R_I(\text{id}_I))$$

$$\eta_x(S) = (S_I(\text{id}_I), S_I(\text{id}_I)).$$

Suppose that $\eta_x(R) = \eta_x(S)$, and let Y be any space, and $\lambda \in QY$.

If λ is a map $I \rightarrow Y$, then the naturality of R implies that the following diagram is commutative:

$$\begin{array}{ccc} Q(I) & \xrightarrow{R_I} & X \wedge I \\ \downarrow Q(\lambda) & & \downarrow X \wedge \lambda \\ Q(Y) & \xrightarrow{R_Y} & X \wedge Y \end{array}$$

i.e. $X \wedge \lambda(R_I(\text{id}_I)) = R_Y(\lambda)$.

If $\eta_x(S) = \eta_x(R)$, we have $R_I(\text{id}_I) = S_I(\text{id}_I)$, so that $R_Y(\lambda) = S_Y(\lambda)$. A similar thing happens if λ is a map $J \rightarrow Y$, so that in all cases, $R_Y(\lambda) = S_Y(\lambda)$, i.e. $R = S$.

Thus η_x is a monomorphism. This and the fact that $\eta_x\sigma_x$ is a homeomorphism imply that both η_x and σ_x are homeomorphisms, i.e. σ and η are natural equivalences.

Hence $T \simeq DQ$, which proves that T , (and hence Q) is reflexion. It is clear that if X is a finite C.W. complex, TX is also a finite C.W. complex (cf. §1).

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