# ON COLIMITS OVER ARBITRARY POSETS 

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#### Abstract

We characterize those partially ordered sets $I$ for which the canonical maps $M_{i} \rightarrow \operatorname{colim} M_{j}$ into colimits of abstract sets are always injective, provided that the transition maps are injective. We also obtain some consequences for colimits of vector spaces.


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1. Introduction. Crowns arise in various problems related to partially ordered sets (posets). Thus, for example, they appear in the study of retracts and fixed points (see [7]), in calculation of the cohomological dimension (see [5]), in applications to homotopy theory (see [11]) and in the investigation of incidence algebras and their quotients (see [1] and [6]). At the same time, quite often they play a "negative" role: the absence of crowns of some kind ensures the existence of certain good properties of posets or constructions related to them. For instance, an incidence algebra $\kappa[S]$ of a finite poset $S$ is completely separating if and only if $S$ contains no crowns [6]. It is not surprising that such a situation arises in a problem of colimits which is discussed in this note: roughly speaking, the crowns are antagonists of directed sets for which colimits are usually considered and well understood (note that colimits over directed posets are called directed colimits, or direct limits, or inductive limits).

More precisely, if one takes a directed colimit colim $M_{i}$ (also denoted by $\underset{\longrightarrow}{\lim } M_{i}$ ), where $i$ runs over a directed poset $I$, such that the transition maps $\varphi_{i j}: M_{i} \rightarrow M_{j}$ are injective, then the canonical maps $M_{i} \rightarrow$ colim $M_{j}$ are also injective, which is a crucial property for applications. Thus, one may wonder which are the posets $I$ for which this always happens. We completely characterize such $I$ in the case when the $M_{i}$ 's are abstract sets and obtain consequences for the colimits of vector spaces.

This problem is related to a similar question about the ring-theoretic version of cross-sectional algebras of Fell bundles over inverse semigroups studied in [8], since such algebras are epimorphic images of colimits of vector spaces over non-necessarily directed posets. In the $C^{*}$-algebraic context, it is proved that the fibres are canonically embedded into the cross-sectional algebra; however, the abstract ring theoretic version
of this fact is not established so far. Moreover, colimits are used in formal software development. This was inspired by J. Goguen in [9] (see also the fifths "dogma" in [10]), and being pursued by several authors (see, in particular, $[4,13,14]$ ). More precisely, colimits are the algebraic tools to compose specifications into new ones, and it is important to make sure that the initial specifications are canonically embedded into more complex ones obtained via general colimits. We believe that the present article may be useful for the above-mentioned problems.
2. Setting of the problem. Let $I$ be a set with a partial order $\unlhd$, considered as a small category, in particular, we interpret $i \unlhd j$ as a (unique) arrow from $i$ to $j$. Denote by Set the category of sets. In analogy with the case of directed posets (see [12, Chapter VIII]), by an $I$-spectrum ${ }^{1} M_{I}$ we mean a diagram $M: I \rightarrow$ Set (i.e. $M$ is a functor) such that each $M(i \unlhd j): M(i) \rightarrow M(j)$ is injective. Write, $M_{i}=M(i)$ and $\varphi_{i j}=M(i \unlhd j)$. The maps $\varphi_{i j}: M_{i} \rightarrow M_{j}$, with $i \unlhd j$, are called the transition maps.

Alternatively, an $I$-spectrum $M_{I}$ consists of (possibly empty) sets $M_{i}(i \in I)$ and maps $\varphi_{i j}: M_{i} \rightarrow M_{j}(i \unlhd j)$ such that

$$
\begin{align*}
& \forall i \in I \quad \varphi_{i i}=\text { id (the identity map), }  \tag{1}\\
& (i \unlhd j \unlhd k) \Longrightarrow \varphi_{j k} \varphi_{i j}=\varphi_{i k}  \tag{2}\\
& \forall i, j \in I(i \unlhd j) \varphi_{i j} \text { isinjective. } \tag{3}
\end{align*}
$$

We shall write $a \triangleleft b$ if $a \unlhd b$ and $a \neq b$, and $a \unlhd b$ will be also denoted by $b \unrhd a$.

REMARK 1. If $M_{i}=\varnothing$ and $i \unlhd j$, then $\varphi_{i j}$ is an empty map and its composition with other maps is also empty. In particular, $\varphi_{i i}$ is the identity map. On the other hand, if $M_{i} \neq \varnothing$ and $i \triangleleft j$, then $M_{j}$ cannot be empty.

Given an $I$-spectrum, the colimit of the diagram $M \rightarrow$ Set will be denoted by colim $M_{i}$. We say that the $I$-spectrum $M_{I}$ is faithful if all canonical maps $M_{i} \rightarrow$ colim $M_{j}$ are injective. A poset $I$ is called faithful if any $I$-spectrum is faithful.

Write $M=\coprod_{i \in I} M_{i}$ and let $\varepsilon_{i}: M_{i} \rightarrow M$ be the natural embeddings. Recall that by the construction of colimits in Set,

$$
\operatorname{colim} M_{i}=M / v,
$$

where $v$ is the smallest equivalence relation containing the following binary relation:

$$
\mu=\left\{\left(\varepsilon_{i}\left(a_{i}\right), \varepsilon_{j} \circ \varphi_{i j}\left(a_{i}\right)\right) \mid i, j \in I, i \unlhd j, a_{i} \in M_{i}\right\} \subseteq M \times M
$$

and the canonical maps $M_{i} \rightarrow \operatorname{colim} M_{j}=M / v$ are the compositions $M_{i} \hookrightarrow M \xrightarrow{\chi}$ $M / \nu$, where $\chi$ is the natural projection. Therefore, $I$ is faithful if and only if for every $I$-spectrum $M_{I}=\left(\left\{M_{i} \mid i \in I\right\},\left\{\varphi_{i j}: M_{i} \rightarrow M_{j} \mid i \unlhd j\right\}\right)$ the maps $\psi_{i}=\chi \circ \varepsilon_{i}$ are all injective.

We shall consider the following:
Problem. Characterize the faithful posets.

[^0]A similar problem can be formulated in the category of vector spaces. Then, we assume that each $M_{i}$ is a vector space over a fixed field and all $\varphi_{i j}$ are linear maps. In this case, we take $M=\oplus_{i \in I} M_{i}$ and the subspace $N \subseteq M$ generated by all $\varepsilon_{i}\left(a_{i}\right)-\varepsilon_{j} \circ \varphi_{i j}\left(a_{i}\right)$, $i, j \in I, i \unlhd j, a_{i} \in M_{i}$. Then, colim $M_{i}=M / N$. We shall say that a poset $I$ is linearly faithful if for any $I$-spectrum of vector spaces each $\psi_{i}$ is injective.

In all what follows, $M_{i}(i \in I)$ will be abstract sets and $\varphi_{i j}: M_{i} \rightarrow M_{j}$ maps of sets.
If $M$ is a set and $\kappa$ a field, denote by $\langle M\rangle$ the vector $\kappa$-space whose base is $M$. Then, $\langle-\rangle$ determines a functor from the category of sets to that of $\kappa$-spaces. The forgetful functor from $\kappa$-spaces to sets is a right adjoint to $\left\langle \_\right\rangle$, and by [3, Corollary 3.9] we have that $\left\langle_{-}\right\rangle$commutes with colimits. Let now $M_{I}=\left(\left\{M_{i} \mid i \in I\right\},\left\{\varphi_{i j} \mid i \unlhd j\right\}\right)$ be an $I$-spectrum, and let $\tilde{\varphi}_{i j}:\left\langle M_{i}\right\rangle \rightarrow\left\langle M_{j}\right\rangle,(i, j \in I, i \unlhd j)$ be the linear extension of $\varphi_{i j}$, i.e. $\tilde{\varphi}_{i j}=\left\langle\varphi_{i j}\right\rangle$. Then, obviously $V_{I}=\left(\left\{\left\langle M_{i}\right\rangle \mid i \in I\right\},\left\{\tilde{\varphi}_{i j} \mid i \unlhd j\right\}\right)$ is an $I$-spectrum of vector spaces, and by the above

$$
\operatorname{colim}\left\langle M_{i}\right\rangle=\left\langle\operatorname{colim} M_{i}\right\rangle .
$$

Consequently, if $M_{I}$ is not faithful, then neither is $V_{I}$, and thus we have the following:

## Proposition 1. If the poset I is linearly faithful, then I is faithful.

Recall that a subset $F$ of a poset $I$ is called a filter (resp. an ideal) if $a \unlhd x$ (resp. if $a \unrhd x$ ) implies $x \in F$ for any $a \in F, x \in I$. If $A \subseteq F$ and for every $x \in F$, there exists $a \in A$, such that $a \unlhd x$, then we say that the filter $F$ is generated by the subset $A$. If $A$ consists of a single element, $A=\{a\}$, then $F$ is called the principal filter generated by $a$.

The set of all common upper bounds of a subset $J \subseteq I$ will be denoted by $U(J)$. If, for example, $J=\{a, b\}$, we also write $U(a, b)$ instead of $U(J)$.
3. Crowns. A poset $K$ is called a crown [7], if $K$ is an union of disjoint sets

$$
\begin{equation*}
K_{0}=\{1,2, \ldots, n\}, \quad K_{1}=\{\overline{1}, \overline{2}, \ldots, \bar{n}\} \tag{4}
\end{equation*}
$$

for some $2 \leq n<\infty$ such that the following relations of comparability, and only they (except the relation of equality), are valid:

$$
\forall \alpha(1 \leq \alpha \leq n) \quad \alpha \triangleleft \bar{\alpha}, \quad \alpha \triangleleft \overline{\alpha+1},
$$

where the indices are taken modulo $n$, i. e. $n+1=1$ and $\overline{n+1}=\overline{1}$. All elements in $K_{0}$ and $K_{1}$ are assumed to be distinct, i.e. $K_{0} \cup K_{1}$ has $2 n$ elements. The diagram of the crown is as follows:


We refer to $n$ as the degree of the crown. Without loss of generality, for any $I$ spectrum $M: I \rightarrow$ Set, the sets $M_{i}$ will be assumed to be pairwise disjoint and the symbol $\varepsilon_{i}$, identifying any $a \in M_{i}$ with its image $\varepsilon_{i}(a) \in \coprod M_{i}$, will be omitted.

The following statement shows the existence of non-faithful posets:

## Proposition 2. No crown is faithful.

Proof. For each $k \in K \backslash\{\overline{1}\}$, we take a singleton $\left\{a_{k}\right\}$ as $M_{k}$. In addition, we set $M_{\overline{1}}=\left\{a_{\overline{1}}, b\right\}$. Thus, the bijections $\varphi_{k l}$ for $(k, l) \neq(1, \overline{1}),(n, \overline{1})$ are determined uniquely. Further, let

$$
\varphi_{1 \overline{1}}\left(a_{1}\right)=a_{\overline{1}}, \quad \varphi_{n \overline{1}}\left(a_{n}\right)=b .
$$

Clearly, the collections $\left\{M_{k}\right\},\left\{\varphi_{k l}\right\}$ form a $K$-spectrum. All pairs

$$
\left(a_{1}, a_{\overline{1}}\right),\left(a_{1}, a_{\overline{2}}\right),\left(a_{2}, a_{\overline{2}}\right),\left(a_{2}, a_{\overline{3}}\right), \ldots,\left(a_{n}, a_{\bar{n}}\right),\left(a_{n}, b\right)
$$

are contained in $\mu$. By the transitivity $\left(a_{\overline{1}}, b\right) \in \nu$. This means that the considered spectrum is non-faithful.

Corollary 1. No crown is linearly faithful.
Proof. This follows using Proposition 1.
However, adding to the crown some additional links may make it faithful, as it is shown in the following statement.

Remark 2. The poset $I$, obtained from the crown $K=\{1,2 ; \overline{1}, \overline{2}\}$ by adding the $\operatorname{link} \overline{1} \triangleleft \overline{2}$, is faithful. This follows from the fact that $I$ is a directed set ([2], Section 7, no 6, Remark 1).
4. Pure subcrowns. Let $I$ be a poset and $J \subseteq I$ a subset. Then, the partial order $\unlhd$ of $I$ induces a partial order $\unlhd_{J}=\unlhd \cap(J \times J)$ on $J$. Then, we shall say that $J$ is a subposet of $I$. A subcrown of $I$ is a subposet of $I$ which is a crown with respect to the induced partial order.

A subcrown $K$ of $I$ of degree $n \geq 4$ will be called pure if

$$
(\forall i, j \in K) \quad U(i, j) \neq \varnothing \Longrightarrow K \cap U(i, j) \neq \varnothing .
$$

It is easy to see that this condition is equivalent to the following:

$$
\begin{equation*}
\left(\forall i, j \in K_{0}\right) \operatorname{dist}(i, j) \geq 4 \Longrightarrow U(i, j)=\varnothing, \tag{5}
\end{equation*}
$$

where $\operatorname{dist}(i, j)$ is the length of the shortest path between $i$ and $j$ in the undirected graph corresponding to the subcrown.

For $n=3$, we shall strengthen the definition of a pure subcrown. Namely, a subcrown $K \subseteq I$ of degree 3 is called pure if $U\left(K_{0}\right)=\varnothing$. We mention by the way that any crown of degree 3 obviously satisfies (5).

For the case $n=2$ we also need a special definition of purity. Let $J$ be a subset of $I$ and $i, k \in J$. We shall say $i$ and $k$ are $J$-connected, if there are $j_{1}, \ldots, j_{m} \in J$, such that

$$
i \stackrel{j_{1}}{\triangleleft} \downarrow \ldots \stackrel{\downarrow}{\triangleleft} j_{m} \stackrel{\triangleright}{\triangleleft},
$$

where every two neighbouring elements are connected by a relation $\unlhd$ or $\unrhd$. If there is no such a sequence, we say that $i$ and $k$ are $J$-disconnected.

A subcrown $K=\{1,2 ; \overline{1}, \overline{2}\} \subseteq I$ of degree 2 is called pure if the elements $\overline{1}$ and $\overline{2}$ are $U(1,2)$-disconnected.

Theorem 1. If I is non-faithful, then I contains a pure subcrown.
Proof. Let $M_{I}=\left(\left\{M_{i}\right\},\left\{\varphi_{i j}\right\}\right)$ be a non-faithful $I$-spectrum.
(1) We show first the existence of a subcrown. There exists an index, say $1 \in I$, and elements $a, b \in M_{1}, a \neq b$ with $\psi_{1}(a)=\psi_{1}(b)$. It follows that $(a, b) \in v$, i. e. there are non-necessarily distinct indices $i_{2}, \ldots, i_{m} \in I$ and elements $c_{2} \in$ $M_{i_{2}}, \ldots, c_{m} \in M_{i_{m}}$ such that

$$
\begin{equation*}
\left(a, c_{2}\right),\left(c_{2}, c_{3}\right), \ldots,\left(c_{m-1}, c_{m}\right),\left(c_{m}, b\right) \in \tilde{\mu} \tag{6}
\end{equation*}
$$

where $\tilde{\mu}$ is the symmetric closure of $\mu$ (i. e. $\tilde{\mu}=\mu \cup \mu^{-1}$ ). It will be convenient to write $2=i_{2}, \ldots, m=i_{m}$ and $i_{1}=i_{m+1}=1, c_{1}=a, c_{m+1}=b$. Then, from (6) we obtain the following chain

$$
\begin{equation*}
1 \triangleright 2 \triangleright \ldots \stackrel{\triangleright}{\triangleleft} m_{\triangleleft}^{\triangleright} 1, \tag{7}
\end{equation*}
$$

and for each $i=2, \ldots, m-1$ one has $c_{i+1}=\varphi_{i, i+1}\left(c_{i}\right)$ or $c_{i}=\varphi_{i+1, i}\left(c_{i+1}\right)$. This will be denoted by $c_{i} \rightarrow c_{i+1}$ or $c_{i} \leftarrow c_{i+1}$, respectively.
We choose $M_{1}$ and $a, b \in M_{1}\left(a \neq b, \psi_{1}(a)=\psi_{1}(b)\right)$ for which the sequence (6) has the smallest length.

First, note that in a minimal chain the signs $\triangleleft$ and $\triangleright$ alternate. Secondly, for the pairs $(1,2)$ and $(m, 1)$ of $(7)$ the signs are different. Indeed, if, for example, $1 \triangleleft 2$ and $m \triangleleft 1$, then we obtain a chain

$$
\left(c_{2}, c_{3}\right), \ldots,\left(c_{m-1}, c_{m}\right),\left(c_{m}, c_{2}^{\prime}\right)
$$

where $c_{2}^{\prime}=\varphi_{12}(b) \neq \varphi_{12}(a)=c_{2}$ (since $\left.a \neq b\right)$. This chain is shorter than (6), a contradiction.
Consequently, the chain of indices (7) has one of the two forms:

$$
1 \triangleleft 2 \triangleright 3 \triangleleft \ldots \triangleleft m \triangleright 1
$$

or

$$
\begin{equation*}
1 \triangleright 2 \triangleleft 3 \triangleright \ldots \triangleright m \triangleleft 1 . \tag{8}
\end{equation*}
$$

In both cases $m$ is even, and it is easy to see that $m \neq 2$.
Observe that without loss of generality, we may assume that our chain is of form (8). Indeed, suppose that the minimal sequence of elements is

$$
\begin{equation*}
a \rightarrow c_{2} \leftarrow c_{3} \rightarrow \ldots \rightarrow c_{m} \leftarrow b . \tag{9}
\end{equation*}
$$

In this case, $1 \triangleleft 2$. Denote $d=\varphi_{12}(b)$. Then, $c_{2} \neq d$ (otherwise, $a=b$ in view of the injectivity) and we obtain a chain of form (8) with the same length as (9):

$$
c_{2} \leftarrow c_{3} \rightarrow \ldots \rightarrow c_{m} \leftarrow b \rightarrow d
$$

A priori, some elements of the sequence (8) may coincide. We shall show that for a minimal chain this is impossible, and, moreover, that the elements in (8) form a subcrown.
Suppose $i \unlhd j$ for some $2 \leq i, j \leq m,|i-j|>1$, i. e. there exists $\varphi_{i j}: M_{i} \rightarrow M_{j}$. Assume $i<j$. Then, if $\varphi_{i j}\left(c_{i}\right)=c_{j}$, we obtain a shorter chain

$$
a \leftarrow c_{2} \rightarrow \ldots \rightleftarrows c_{i} \rightarrow c_{j} \rightleftarrows \ldots \leftarrow c_{m} \rightarrow b
$$

(here the symbol $\rightleftarrows$ denotes one of the two arrows $\leftarrow$ or $\rightarrow$ ). If $\varphi_{i j}\left(c_{i}\right) \neq c_{j}$, then $\left(c_{i}, c_{j}\right) \in \nu$ and again we get a shorter chain

$$
\varphi_{i j}\left(c_{i}\right) \leftarrow c_{i} \rightleftarrows \ldots \rightleftarrows c_{j}
$$

The case $i>j$ is symmetric.
All these arguments remain valid if we take $a$ instead of $c_{i}$ (or of $c_{j}$ ), noting that in this case $i=1$ and $3 \leq j \leq m-1$ ( or $j=1$ and $3 \leq i \leq m-1$ ). This shows that (8) gives a subcrown.
(2) We show next that the obtained subcrown $K$ is pure for $m \geq 8$ (i. e. when the degree of the crown $\geq 4$ ). It suffices to verify the condition (5). Take arbitrary $i, j \in K_{0}$ with $\operatorname{dist}(i, j) \geq 4$ and suppose that there exists $k \in I$ such that $i \unlhd k$, $j \unlhd k$. Assume $i<j$. Then, we have

$$
\ldots \rightarrow c_{i-1} \leftarrow c_{i} \rightarrow c_{i+1} \leftarrow \ldots \rightarrow c_{j-1} \leftarrow c_{j} \rightarrow c_{j+1} \leftarrow \ldots
$$

If $\varphi_{i k}\left(c_{i}\right)=\varphi_{j k}\left(c_{j}\right)=d$, then the original chain can be shortened:

$$
\ldots \rightarrow c_{i-1} \leftarrow c_{i} \rightarrow d \leftarrow c_{j} \rightarrow c_{j+1} \leftarrow \ldots
$$

If $\varphi_{i k}\left(c_{i}\right) \neq \varphi_{j k}\left(c_{j}\right)$, we get a new shorter chain:

$$
\varphi_{i k}\left(c_{i}\right) \leftarrow c_{i} \rightarrow c_{i+1} \leftarrow \ldots \rightarrow c_{j-1} \leftarrow c_{j} \rightarrow \varphi_{j k}\left(c_{j}\right)
$$

The case $i>j$ is similar.
(3) We check now the purity of $K$ for $m=6$. Consider the chain (8). Suppose that there is an element $i \in I$ such that $2 \triangleleft i, 4 \triangleleft i, 6 \triangleleft i$. Write

$$
\varphi_{2 i}\left(c_{2}\right)=x, \quad \varphi_{4 i}\left(c_{4}\right)=y, \quad \varphi_{6 i}\left(c_{6}\right)=z .
$$

If $x \neq y$, we obtain a shorter chain

$$
x \leftarrow c_{2} \rightarrow c_{3} \leftarrow c_{4} \rightarrow y .
$$

Therefore, $x=y$ and, similarly, $y=z$. But then we get again the shorter chain

$$
a \leftarrow c_{2} \rightarrow x \leftarrow c_{6} \rightarrow b,
$$

contradicting our assumption.
(4) It remains to consider the case $m=4$. Suppose that our crown $K=\{2,4 ; 1,3\}$ has form (8):

$$
1 \triangleright 2 \triangleleft 3 \triangleright 4 \triangleleft 1
$$

with the corresponding sequence of elements

$$
a \leftarrow c_{2} \rightarrow c_{3} \leftarrow c_{4} \rightarrow b,
$$

where $a \neq b$.
Assume that the elements 1 and 3 are $U(2,4)$-connected. Write $j_{0}=1$ and $j_{l+1}=3$, and let $j_{1}, \ldots, j_{l} \in U(2,4)$ be such that

$$
j_{0} \triangleright j_{1} \triangleright \ldots{ }^{\triangleright} j_{l} \triangleright j_{l} \triangleright j_{l+1} .
$$

Consider the "intermediate" subposets $K^{\alpha}=\left\{2,4 ; j_{\alpha}, j_{\alpha+1}\right\}, 0 \leq \alpha \leq l$. By Remark 2, they are faithful. Therefore, $\left.\left(\varphi_{2 j_{l}}\left(c_{2}\right), \varphi_{4 j l}\left(c_{4}\right)\right) \in \nu\right|_{K^{\prime}}$ implies $\varphi_{2 j_{l}}\left(c_{2}\right)=$ $\varphi_{4 j l}\left(c_{4}\right)$. The latter equality, in turn, implies $\varphi_{2 j_{l-1}}\left(c_{2}\right)=\varphi_{4 j-1}\left(c_{4}\right)$ and so on. At the last step, we get $a=\varphi_{21}\left(c_{2}\right)=\varphi_{41}\left(c_{4}\right)=b$, contradicting the assumption $a \neq b$.
5. Filters and faithfulness. For the converse of Theorem 1, we need additional notions.

Let $I$ be a poset. Denote by $\min I$ the set of minimal elements of $I$. For $a \in I$, let $F_{a}=\{x \in I \mid x \unrhd a\}$ (the principal filter generated by $a$ ).

Consider the following conditions on $I$ :
(1) $|\min I|<\infty$. In this case, we shall use positive integers to denote the minimal elements, i. e. we write $\min I=\{1,2, \ldots, n\}$.
(2) $\bigcup_{i \in \min I} F_{i}=I$.
(3) For all $i, j \in \min I, i \neq j$,

$$
F_{i} \cap F_{j} \begin{cases}\neq \varnothing, & \text { if }|i-j| \leq 1 \\ =\varnothing & \text { otherwise }\end{cases}
$$

(as in Section 3 we work with $i, j$ modulo $n$ ). In particular, if $n \geq 4$ then $F_{i} \cap$ $F_{j} \cap F_{k}=\varnothing$ for pairwise different $i, j, k \in \min I$.
For $n=3$, the last equality may not hold, so we add one more condition:
(4) If $n=3$, then $F_{1} \cap F_{2} \cap F_{3}=\varnothing$.

Proposition 3. If a poset I satisfies conditions (1)-(4) and $n \geq 3$, then I is nonfaithful.

Proof. We construct an $I$-spectrum $M_{I}=\left(\left\{M_{x} \mid x \in I\right\},\left\{\varphi_{x y} \mid x \unlhd y\right\}\right)$ as follows.
Write $G_{i}=F_{i-1} \cap F_{i}$ (in particular, $G_{1}=F_{n} \cap F_{1}$ ). Note that by conditions (3) and (4), $G_{i} \cap G_{j} \neq \varnothing$ implies $i=j$. We put

$$
M_{x}= \begin{cases}\left\{a_{x}\right\} & \text { if } x \notin G_{1} \\ \left\{a_{x}, b_{x}\right\} & \text { if } x \in G_{1}\end{cases}
$$

We assume that $a_{x} \neq b_{x}$ and that the sets $M_{x}$ are disjoint.
The maps $\varphi_{x y}(x \unlhd y)$ are defined uniquely, if $\left|M_{y}\right|=1$. For $\left|M_{x}\right|=\left|M_{y}\right|=2$, $x \unlhd y$, we set

$$
\begin{equation*}
\varphi_{x y}\left(a_{x}\right)=a_{y}, \quad \varphi_{x y}\left(b_{x}\right)=b_{y} \tag{10}
\end{equation*}
$$

Next, we define

$$
\begin{equation*}
\varphi_{x z}\left(a_{x}\right)=a_{z}, \quad \varphi_{y z}\left(a_{y}\right)=b_{z} \tag{11}
\end{equation*}
$$

for any $x \in F_{1} \backslash G_{1}, y \in F_{n} \backslash G_{1}, z \in G_{1}$ such that $x \triangleleft z$ and $y \triangleleft z$. Since $y \unrhd x \in G_{1}$ implies $y \in G_{1}$, it follows by conditions (2)-(4) that all $\varphi_{x y}$ are completely defined. Evidently, each $\varphi_{x y}$ is injective.

To prove that $M_{I}$ is an $I$-spectrum, we have to verify the condition

$$
\begin{equation*}
(x \unlhd y \unlhd z) \Longrightarrow \varphi_{y z} \varphi_{x y}=\varphi_{x z} . \tag{12}
\end{equation*}
$$

If either $x, y, z \notin G_{1}$ or $x, y, z \in G_{1}$, then (12) is evident. Therefore, we can assume that $z \in G_{1}$ (otherwise $x, y \notin G_{1}$ ) and $x \notin G_{1}$ (otherwise $y, z \in G_{1}$ ). Then, from condition (3) we have $x \in F_{n} \backslash G_{1}$ or $x \in F_{1} \backslash G_{1}$. In the first case, one have either $y \in F_{n} \backslash G_{1}$ or $y \in G_{1}$, and (12) is true. The second case is treated the same way.

Choosing an arbitrary element $\bar{k} \in G_{k}$ for each $k \leq n$, we get a crown

$$
K=\{1, \ldots, n ; \overline{1}, \ldots, \bar{n}\} .
$$

Now the non-faithfulness of $M_{I}$ follows from the fact that its restriction to $K$ coincides with the spectrum constructed in the proof of Proposition 2.
6. The converse of Theorem 1. We begin with the following statement.

Lemma 1. Each filter in a faithful poset is faithful.
Proof. Let $F$ be a filter in a faithful poset $I$, and let

$$
M_{F}=\left(\left\{M_{x} \mid x \in F\right\},\left\{\varphi_{x y} \mid x \unlhd y\right\}\right)
$$

be an $F$-spectrum. We extend it to $I$ by putting $M_{x}=\varnothing$ for $x \in I \backslash F$. The maps $\varphi_{x y}$ for $x \in I \backslash F, x \unlhd y$, are determined automatically and it is easy to see that we obtained an $I$-spectrum $M_{I}$ which must be faithful. Therefore, $M_{F}$ is also faithful.

## Theorem 2. If a poset I contains a pure subcrown, then I is non-faithful.

Proof. Let $K$ be a pure subcrown in $I$ of form (4) (see Section 3). By Lemma 1, replacing $I$ by the filter generated by $K_{0}=\{1,2, \ldots, n\}$, we may assume that $\bigcup_{i \in K_{0}} F_{i}=$ $I$. In particular, $\min I=K_{0}$.

First, suppose that $n \geq 3$. We have that $I$ satisfies conditions (1) and (2) from Section 5. Furthermore, $G_{k} \neq \varnothing$ for any $k \in K_{0}$, since it contains the element $\bar{k}$ of the crown. If $F_{i} \cap F_{j} \neq \varnothing$, then it follows from the purity that $|i-j| \leq 1$, i. e. (3) holds. Finally, condition (4) follows from the definition of purity for $n=3$. Then by Proposition 3, we conclude that $I$ is non-faithful.

It remains to consider the case $n=2$. Let $K=\{1,2 ; \overline{1}, \overline{2}\}$ be a pure subcrown in $I$ and $I=F_{1} \cup F_{2}$, where $F_{1}$ and $F_{2}$ are the principal filters generated by 1 and 2. Recall that $\min I=\{1,2\}$. Write $G=F_{1} \cap F_{2}=U(1,2)$. Denote by $H$ the set of elements of $G$ which are $G$-connected with $\overline{1}$. We note some obvious properties of $H$ :
$H$ is a filter in $I$;
$H$ is an ideal in $G$;
$\overline{2} \notin H$, as $K$ is pure.
Similarly as we did in the proof of Proposition 3, set

$$
M_{x}= \begin{cases}\left\{a_{x}\right\}, & \text { if } x \notin H, \\ \left\{a_{x}, b_{x}\right\}, & \text { if } x \in H .\end{cases}
$$

The maps $\varphi_{x y}(x \unlhd y)$ are uniquely defined, if $y \notin H$. For $x, y \in H, x \unlhd y$, they are given by the equalities (10). For $x \in F_{1} \backslash G, y \in F_{2} \backslash G, z \in H$ such that $x \triangleleft z$ and $y \triangleleft z$ the maps $\varphi_{x z}$ and $\varphi_{y z}$ are defined by (11). Since $H$ is a filter in $I$ and an ideal in $G$, all $\varphi_{x y}$ are defined and they are injective.

We need to check (12). If either $x, y, z \notin H$ or $x, y, z \in H$, then (12) is obvious. Therefore, we can assume that $z \in H$ (otherwise $x, y \notin H$ ) and $x \notin H$ (otherwise $y, z \in$ $H)$. Further, if $y \in G$, then $y \in H$ as $H$ is an ideal in $G$, and (12) is readily verified. Assume $y \notin G$.

Suppose $y \in F_{1} \backslash G$. Then, $x \notin F_{2}$, as $y \notin G$. Hence, $x \in F_{1} \backslash G$ and (12) is true. The case $y \in F_{2} \backslash G$ is symmetric. It remains, as in the Proposition 3, to refer to the proof of Proposition 2.

Applying Proposition 1, we obtain the following:
Corollary 2. If a poset I contains a pure subcrown, then I is not linearly faithful.

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[^0]:    ${ }^{1}$ One may also use the term inductive system from Category Theory.

