DECOMPOSITION OF GRAPHS INTO TWO-WAY INFINITE PATHS

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1. Introduction. We consider undirected graphs in which two vertices may be joined by more than one edge and in which a vertex may be joined to itself by one or more edges. G will always denote a graph. The set of vertices [edges] of G will be denoted by V(G) [E(G)]. G is finite or infinite according as $V(G) \cup E(G)$ is finite or infinite. The degree, $d(\xi)$ or $d_{G}(\xi)$, of a vertex ξ of G is 2a + b, where a is the number of edges joining ξ to itself and b is the number of those joining ξ to other vertices. A vertex is *even* [odd] if it has even [odd] finite degree. A subgraph of G is a graph H such that $V(H) \subset V(G)$, $E(H) \subset E(G)$, and each edge of H joins the same vertices in H as in G. A *component* of G is a maximal connected subgraph. A collection of subgraphs of G are disjoint [edge-disjoint] if no two have a common vertex [edge]. |A|denotes the cardinal number of the set A. G will be called relevant if $|E(G)| = \aleph_0$ and G has no odd vertex or finite component. A decomposition of G is a set of edge-disjoint subgraphs of G whose union is G. A *path-sequence* of G is a finite or infinite sequence whose terms are alternately vertices and edges of G, starting and ending if at all with a vertex, such that each edge appearing in the sequence appears just once and joins the vertices immediately preceding and following it. An infinite path-sequence is one-way infinite [two-way infinite] if it has [has not] a first or last term. The subgraph formed by the terms of a path-sequence [*n*-way infinite path-sequence] is a *path* [*n*-way *infinite path* (n = 1, 2)]. (The same path can be both one-way and two-way infinite.) It follows from (1, Theorem 2) that a relevant graph is decomposable into two-way infinite paths, and Ore (3, p. 47, problem 3) asks in effect what is the minimum number of such paths into which it can be decomposed. This paper answers the question.

2. Limited graphs. DEFINITIONS. If X, Y are subsets of V(G), \bar{X} will denote V(G) - X, $X \circ Y$ will denote the set of those edges of G which join elements of X to elements of Y, $X\delta$ (or $X\delta_G$) will denote $X \circ \bar{X}$, and X^* will denote the subgraph of G defined by $V(X^*) = X$, $E(X^*) = X \circ X$. X is *inessential* if $X \circ V(G)$ is finite. $X\delta$ will be called a *cincture*; it is an *even* [odd] cincture if $|X\delta|$ is finite and even [odd]. Whenever two or more graphs are under consideration and one of them is denoted by G, the notations $d(\xi)$, $X\delta$ will mean $d_G(\xi)$, $X\delta_G$, and the notations $X \circ Y$, \bar{X} , X^* will also be interpreted relative to G. (But an odd vertex of a subgraph of G is one whose degree in

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that subgraph is odd.) If A_1, \ldots, A_n are sets, $A_1 + \ldots + A_n$ denotes the set of those elements which belong to an odd number of the A_i . Two subgraphs H, K of G are similar (in symbols, $H \sim K$) if the symmetric differences V(H) + V(K) and E(H) + E(K) are finite. Clearly similarity is an equivalence relation and so subdivides the subgraphs of G into similarity classes. Let l be a non-negative integer. An *l-splitting* of G is a set $\{H_1, \ldots, H_l\}$ of disjoint infinite subgraphs of G such that G is the union of these and a finite subgraph: any finite subgraph H such that $G = H_1 \cup \ldots \cup H_l \cup H$ will be called a *completion* of the splitting. G is *l-limited* (in symbols, $\lim G = l$), if it has an *l*-splitting but no l'-splitting for l' > l. G is *limited* if it is *l*-limited for some non-negative integer l.

LEMMA 1. If G is l-limited, there are l distinct similarity classes of subgraphs of G such that the subgraphs in any l-splitting of G belong one to each of these classes.

Proof. Since the members of an *l*-splitting of *G* are clearly dissimilar, it suffices to prove that, if $\{H_1, \ldots, H_l\}$ and $\{K_1, \ldots, K_l\}$ are *l*-splittings of *G*, each H_i is similar to one of the K_j . To prove this, let the given splittings have completions *H*, *K* respectively. Since the subgraphs $H_i \cap K_j$ are disjoint and the union of these and the finite subgraph $H \cup K$ is *G*, which is *l*-limited, at most *l* of the $H_i \cap K_j$ can be infinite. But, since the H_i and K_j are infinite. Hence there is a permutation K_1', \ldots, K_l' of K_1, \ldots, K_l such that $H_i \cap K_j$ is infinite if and only if i = j. Since $V(H_i) + V(K_i')$, $E(H_i) + E(K_i')$ are contained in the sets of vertices and edges of the finite graph

$$H \cup K \cup \bigcup_{j \neq i} [(H_i \cap K_j') \cup (H_j \cap K_i')],$$

it follows that $H_i \sim K_i'$ (i = 1, ..., l) as required.

DEFINITIONS. The l similarity classes characterized by Lemma 1 will be called wings of G. If W is a wing of G and X is a subset of V(G) such that $X^* \in W$ and X δ is finite, we shall call X a W-set and X δ a W-cincture.

LEMMA 2. If X is an inessential subset of V(G), then

$$|X\delta| \equiv \sum_{\xi \in X} d(\xi) \pmod{2}.$$

Proof. This is so since an edge contributes 2, 1, or 0 to $\sum_{\xi \in X} d(\xi)$ according as it belongs to $X \circ X$, $X\delta$, or $\overline{X} \circ \overline{X}$, respectively.

COROLLARY 2A. If, in addition, G is relevant, $X\delta$ is even.

LEMMA 3. If X, $Y \subset V(G)$, then $X\delta + Y\delta = (X + Y)\delta$.

The proof is left to the reader.

LEMMA 4. Let G be l-limited, $\{H_1, \ldots, H_l\}$ be an l-splitting of G, and W_i be

the wing of G which includes H_i . Then (i) H_i is 1-limited and (ii) $V(H_i)$ is a W_i -set.

Proof. (i) Obviously H_i has a 1-splitting. On the other hand, the elements of any *r*-splitting of H_i together with the H_j $(j \neq i)$ constitute an (l + r - 1)-splitting of G; hence r cannot exceed 1.

(ii) If H is a completion of the given splitting, $E(V(H_i)^*) - E(H_i)$ and $V(H_i)\delta$ are both contained in the finite set E(H). Hence $V(H_i)^* \sim H_i \in W_i$ and $V(H_i)$ is a W_i -set.

LEMMA 5. If G is limited and relevant and W is a wing of G, there exists a W-cincture and all W-cinctures have the same parity.

Proof. The first assertion follows from Lemma 4(ii). Moreover, if X, Y are W-sets, write $X \cap Y = Z_1$, $X \cap \overline{Y} = Z_2$, $\overline{X} \cap Y = Z_3$, $\overline{X} \cap \overline{Y} = Z_4$, and $Z_i \circ Z_j = A_{ij}$. Then

$$(X + Y) \circ V(G) = \bigcup_{i=1}^{4} (A_{i2} \cup A_{i3}) \subset A_{22} \cup A_{33} \cup \bigcup_{i \neq j} A_{ij},$$

which is finite since $\bigcup_{i\neq j} A_{ij} = X\delta \cup Y\delta$, $A_{22} \cup A_{33} \subset E(X^*) + E(Y^*)$, and $X^* \sim Y^*$. Hence X + Y is inessential and so, by Lemma 3 and Corollary 2A,

$$|X\delta| + |Y\delta| \equiv |X\delta + Y\delta| = |(X + Y)\delta| \equiv 0 \pmod{2}.$$

Therefore $X\delta$ and $Y\delta$ have the same parity, and Lemma 5 is proved.

DEFINITIONS. If G is limited and relevant, a wing W of G will be called odd or even according as the W-cinctures of G are odd or even respectively, and p(G)will denote $\alpha + \frac{1}{2}\beta$, where α is the number of even wings of G and β is the number of odd ones. For an unlimited relevant G, we define p(G) to be \aleph_0 .

LEMMA 6. If G is limited and relevant, p(G) is an integer.

Proof. Let $\lim G = l$ and $\{H_1, \ldots, H_l\}$ be an *l*-splitting of G with completion H. Let W_i be the wing of G which includes H_i and let $X_i = V(H_i)$. Let $Y = V(G) - X_1 \cup \ldots \cup X_l$. By Lemma 3 and induction,

(1)
$$X_1\delta + \ldots + X_l\delta = (X_1 + \ldots + X_l)\delta = (X_1 \cup \ldots \cup X_l)\delta = Y\delta.$$

Since $Y \circ V(G) \subset E(H)$, Y is inessential and hence $Y\delta$ is even by Corollary 2A. Moreover, $X_i\delta$ is a W_i -cincture by Lemma 4(ii). Therefore, by (1), an even number of the W_i are odd and p(G) is an integer.

3. Solution of Ore's problem. Definitions. If H is a subgraph of G, G - H will denote the subgraph obtained by omitting from G the edges of H and those vertices of H which are not incident with other edges. The path formed by the terms of a path-sequence will be said to be *derived* from it. When a path-sequence is denoted by a small letter, the corresponding capital letter will denote the path derived from it. For any edge λ appearing in a two-way infinite path-sequence *s*, the sub-sequence formed by the terms before [after] λ will be called a *tail* [*head*] of *s*. An *open* [*closed*] path-sequence is a finite path-sequence whose first and last terms are different [the same]. An *open* [*closed*] path is one derivable from an open [closed] path-sequence. Clearly an open path has just two odd vertices which are the end terms of every path-sequence from which it is derivable: the path will be said to *connect* these vertices. A $\xi \infty$ -*path* is one derivable from a one-way infinite path-sequence with first or last term ξ .

LEMMA 7. Let $\{H_1, \ldots, H_i\}$ be an *l*-splitting of G and s_1, \ldots, s_k be two-way infinite path-sequences of G such that $G = S_1 \cup \ldots \cup S_k$. Then we can select a tail t_i and head u_i of s_i for $i = 1, \ldots, k$ such that each T_i and U_i is contained in one of the H_j and each H_j contains at least one T_i or U_i .

Proof. Let H be a completion of $\{H_1, \ldots, H_i\}$. Since E(H) is finite, s_i has a tail t_i and head u_i neither of which includes an edge of H. Clearly S_i is the union of T_i , U_i , and a finite subgraph F_i . Since the T_i and U_i are connected infinite graphs having no edges in common with H, each of them is contained in an H_j . It follows, since no H_j could be contained in the finite graph $F_1 \cup \ldots \cup F_k$, that each H_j contains at least one T_i or U_i .

LEMMA 8. If G is decomposable into k two-way infinite paths, where k is a positive integer, then G is limited and relevant and $k \ge p(G)$.

Proof. G is obviously relevant. Let s_1, \ldots, s_k be two-way infinite pathsequences such that the S_i decompose G. If \mathfrak{H} is a splitting of G, then by Lemma 7 there exist 2k non-empty subgraphs such that each member of \mathfrak{H} contains at least one of them, whence $|\mathfrak{H}| \leq 2k$. Therefore G is limited and $\lim G \leq 2k$. Let $\lim G = l$ and $\{H_1, \ldots, H_l\}$ be an *l*-splitting of G. Let W_j be the wing of G which includes H_j and let $X_j = V(H_j)$. By Lemma 7, we can select for $i = 1, \ldots, k$ a tail t_i and head u_i of s_i such that each T_i and U_i is contained in one of the H_i and each H_i contains at least one T_i or U_i . If n_i is the number of the subgraphs $T_1, \ldots, T_k, U_1, \ldots, U_k$ which are contained in H_j , then $n_j \ge 1$. Moreover, if $n_j = 1$, then, for some *i*, one of T_i , U_i is contained in H_i whilst the other, together with all the T_h , U_h $(h \neq i)$, is contained in $\bigcup_{m\neq j} H_m$. Hence one of t_i , u_i has all its vertex-terms in X_j whilst the other, and the t_h , u_h $(h \neq i)$, have them all in \bar{X}_j . Therefore s_i includes an odd number of edges from $X_i \delta$ and each s_h $(h \neq i)$ includes an even number. Therefore $X_i\delta$ is odd, and hence, by Lemma 4(ii), W_j is odd. We have thus seen that $n_i \ge 1$ and can only be 1 if W_i is odd, from which it follows that $p(G) \leq \frac{1}{2} \sum n_j = k.$

LEMMA 9. If G is finite and has k odd vertices, then k is even and G is decomposable into $\frac{1}{2}k$ open paths connecting all the odd vertices in pairs and a finite number of disjoint closed paths.

Proof. By (3, Theorem 3.1.1), each component of G which has no odd

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vertices is a closed path. Moreover, if a component of G has h (>0) odd vertices, then h is even by Lemma 2 or (3, Theorem 1.2.1), and it follows from the proof (although not, strictly speaking, from the statement) of (3, Theorem 3.1.2) that the component is decomposable into $\frac{1}{2}h$ open paths connecting all its odd vertices in pairs. Lemma 9 follows from these remarks.

LEMMA 10. If G is 1-limited and connected and has no odd vertex and $|E(G)| = \aleph_0$, then G is a two-way infinite path.

Proof. This follows from (3, Theorem 3.2.2) if we show that, for any finite subgraph H of G, G - H has just one infinite component. But this is clear, since G - H is infinite but (since G is 1-limited) not the union of two disjoint infinite subgraphs.

LEMMA 11. If $|E(G)| = \aleph_0$, $\xi \in V(G)$, $d(\xi)$ is odd or infinite, and G is 1-limited and connected and has no odd vertex except possibly ξ , then G is a ξ^{∞} -path.

Proof. This follows from the proof of (3, Theorem 3.2.1), since the above proof of Lemma 10 shows that G satisfies Condition δ_1 of that theorem.

LEMMA 12. If G has no finite component and F_0 is a finite subgraph of G, then F_0 is contained in a finite subgraph H such that no component of G - His finite or includes more than one odd vertex of H.

Proof. If a component of $G - F_0$ includes two odd vertices of F_0 , let F_1 be the union of F_0 and a path connecting two such vertices in $G - F_0$. If a component of $G - F_1$ includes two odd vertices of F_1 , let F_2 be the union of F_1 and a path connecting two such vertices in $G - F_1$; and so on. Since F_{i+1} has two less odd vertices than F_i , we must ultimately reach an F_n such that no component of $G - F_n$ includes more than one odd vertex of F_n . If H is the union of F_n and the finite components of $G - F_n$, then no component of G - H is finite or includes more than one odd vertex of H. To see that H is finite, we must show that $G - F_n$ has only finitely many finite components. But this is clear, since each finite component of $G - F_n$, not being a component of G, must include a vertex of F_n .

LEMMA 13. A relevant limited graph G is decomposable into p(G) two-way infinite paths.

Proof. Let $\lim G = l$. Let $\{H, \ldots, H_l\}$ be an *l*-splitting of *G* with completion *H*. By Lemma 12, *H* is contained in a finite subgraph *K* such that no component of G - K is finite or includes more than one odd vertex of *K*. Moreover, G - K is the union of the disjoint infinite subgraphs $(G - K) \cap H_i$, and cannot have more than *l* infinite components since $\lim G = l$. Hence the components of G - K are precisely the infinite subgraphs I_1, \ldots, I_l , where $I_i = (G - K) \cap H_i$. Moreover, $\{I_1, \ldots, I_l\}$ is an *l*-splitting of *G*, and, if W_i denotes the wing of *G* which includes I_i , then I_i is 1-limited and $V(I_i)$

is a W_i -set by Lemma 4. But $V(I_i)\delta = V(I_i \cap K)\delta_K$, which by Lemma 2 is even or odd according as I_i includes 0 or 1 odd vertices of K respectively. Hence I_i includes 0 or 1 odd vertices of K according as W_i is even or odd. In case W_i is odd, let ξ_i be the odd vertex of K in I_i . Since the degree in I_i of a vertex ξ is $d(\xi)$ if $\xi \notin V(K)$ and $d(\xi) - d_K(\xi)$ if $\xi \in V(K)$, I_i has no odd vertex except possibly, if W_i is odd, ξ_i , whose degree in I_i is then odd or infinite. We have seen that I_i is 1-limited. Hence, by Lemmas 10 and 11, I_i is a two-way infinite path if W_i is even and is a $\xi_i \infty$ -path if W_i is odd.

Since a vertex of K which is not in one of the I_i has the same degree in Kas in G, the odd vertices of K are precisely the ξ_i associated with the odd W_i . Therefore, by Lemma 9, K is decomposable into some open paths P_1, \ldots, P_m connecting all these ξ_i in pairs and some disjoint closed paths C_1, \ldots, C_n , where m is half the number of ξ_i , i.e. of odd wings of G. If P_r connects ξ_h to ξ_i , let $P_r' = I_h \cup P_r \cup I_i$, which is a two-way infinite path since I_h is a $\xi_h \infty$ -path and I_i a $\xi_i \infty$ -path. Hence $C_1, \ldots, C_n, P_1', \ldots, P_m'$, and the I_i associated with the even W_i constitute a decomposition of G into p(G) twoway infinite paths and n disjoint closed paths. Each of the closed paths, not being a component of G, has a vertex in common with at least one of the two-way infinite paths. Hence, by absorbing each closed path into an infinite path with which it has a vertex in common, we obtain ultimately a decomposition of G into p(G) two-way infinite paths.

By Lemmas 8 and 13, p(G) is the minimum number of two-way infinite paths into which a limited relevant graph G can be decomposed. Moreover, if G is unlimited and relevant, p(G) still has this property, since then (i) $p(G) = \aleph_0$ by definition, (ii) G is decomposable into two-way infinite paths by Theorem 2 of (1), and (iii) the number of paths in such a decomposition must be \aleph_0 by Lemma 8.

4. The maximum number of paths. In this section, G denotes a relevant graph. We shall call G k-constricted if all vertices have finite degree and V(G) can be partitioned into an infinite sequence of disjoint finite subsets X_1, X_2, \ldots such that $|X_i \circ X_{i+1}| = k$ $(i = 1, 2, \ldots)$ and $X_r \circ X_s = \emptyset$ whenever $|r - s| \ge 2$. If G is k-constricted for some positive integer k, the least such k is its width w(G); if not, $w(G) = \aleph_0$. It is not very hard to show that w(G) is even or \aleph_0 and that $(\inf \frac{1}{2}\aleph_0 \mod \aleph_0)$ G cannot be decomposed into more than $\frac{1}{2}w(G)$ two-way infinite paths. But it was shown in (2) that, for every α such that $p(G) \le \alpha \le \frac{1}{2}w(G)$, there exists a decomposition of G into α two-way infinite paths. A somewhat similar result concerning decompositions of undirected graphs into two-way infinite paths. As, however, the details of all this are very lengthy and perhaps rather tedious, I have here confined myself to answering Ore's question.

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