Dense Complete Set For NP

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Abstract

A sparse language is a formal language such that the number of strings of length n is bounded by a polynomial function of n. We create a class with the opposite definition, that is a class of languages that are dense instead of sparse. We define a dense language on m as a formal language (a set of binary strings) where there exists a positive integer n_0 such that the counting of the number of strings of length $n \geq n_0$ in the language is greater than or equal to 2^{n-m} where m is a real number and $0 < m \leq 1$. We call the complexity class of all dense languages on m as $DENSE(m)$. We prove that there exists an *NP–complete* problem that belongs to $DENSE(m)$ for every possible value of $0 < m \leq 1$.

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1 Summary

In computational complexity theory, a sparse language is a formal language (a set of strings) such that the complexity function, counting the number of strings of length n in the language, is bounded by a polynomial function of n. The complexity class of all sparse languages is called $SPARSE$. $SPARSE$ contains $TALLY$, the class of unary languages, since these have at most one string of any one length.

Fortune showed in 1979 that if any sparse language is *coNP–complete*, then $P = NP$ (this is Fortune's theorem) [5]. Mahaney used this to show in 1982 that if any sparse language is *NP–complete*, then $P = NP$ [6]. A simpler proof of this based on left-sets was given by Ogihara and Watanabe in 1991 [7]. Mahaney's argument does not actually require the sparse language to be in NP, so there is a sparse *NP–hard* set if and only if $P = NP$ [6].

We create a class with the opposite definition, that is a class of languages that are dense instead of sparse. We show there is a sequence of languages that are in *NP–complete*, but their density grows as much as we go forward into the iteration of the sequence. The first element of the sequence is a variation of the *NP–complete* problem known as *HAM–CYCLE* [8]. The next element in the sequence is constructed from this new version of *HAM–CYCLE*. Indeed, each language is created from its previous one in the sequence. Since the density grows according we move forward into the sequence, then there exists a language so much dense such that its density tends to 0 when the bit-length n of the binary strings tends to infinity. However, this incredible dense language is still *NP–complete*.

2 Basic Definitions

Let Σ be a finite alphabet with at least two elements, and let Σ^* be the set of finite strings over Σ [1]. A Turing machine M has an associated input alphabet Σ [1]. For each string w in Σ^* there is a computation associated with M on input w [1]. We say that M accepts w if this computation terminates in the accepting state, that is $M(w) =$ "*yes*" [1]. Note that M fails to accept w either if this computation ends in the rejecting state, that is $M(w) =$ "*no*", or if the computation fails to terminate [1].

The language accepted by a Turing machine M, denoted $L(M)$, has an associated alphabet Σ and is defined by

$$L(M) = \{w \in \Sigma^* : M(w) = \text{“}yes\text{”}\}.$$

We denote by $t_M(w)$ the number of steps in the computation of M on input w [1]. For $n \in \mathbb{N}$ we denote by $T_M(n)$ the worst case run time of M; that is

$$T_M(n) = max\{t_M(w) : w \in \Sigma^n\}$$

where Σ^n is the set of all strings over Σ of length n [1]. We say that M runs in polynomial time if there is a constant k such that for all n, $T_M(n) \leq n^k + k$ [1]. In other words, this means the language $L(M)$ can be accepted by the Turing machine M in polynomial time. Therefore, P is the complexity class of languages that can be accepted in polynomial time by deterministic Turing machines [4]. A verifier for a language L is a deterministic Turing machine M, where

$$L = \{w : M(w, c) = \text{“}yes\text{”} \textit{ for some string } c\}.$$

We measure the time of a verifier only in terms of the length of w, so a polynomial time verifier runs in polynomial time in the length of w [1]. A verifier uses additional information, represented by the symbol c, to verify that a string w is a member of L. This information is called certificate. NP is also the complexity class of languages defined by polynomial time verifiers [8]. If NP is the class of problems that have succinct certificates, then the complexity class $coNP$ must contain those problems that have succinct disqualifications [8]. That is, a “*no*” instance of a problem in $coNP$ possesses a short proof of its being a “*no*” instance [8].

A function $f : \Sigma^* \to \Sigma^*$ is a polynomial time computable function if some deterministic Turing machine M, on every input w, halts in polynomial time with just $f(w)$ on its tape [9]. Let $\{0, 1\}^*$ be the infinite set of binary strings, we say that a language $L_1 \subseteq \{0, 1\}^*$ is polynomial time reducible to a language $L_2 \subseteq \{0, 1\}^*$, written $L_1 \leq_p L_2$, if there is a polynomial time computable function $f : \{0, 1\}^* \to \{0, 1\}^*$ such that for all $x \in \{0, 1\}^*$,

$$x \in L_1 \textit{ if and only if } f(x) \in L_2.$$

An important complexity class is *NP–complete* [4]. A language $L \subseteq \{0, 1\}^*$ is *NP–complete* if

- $L \in NP$, and
- $L' \leq_p L$ for every $L' \in NP$.

If L is a language such that $L' \leq_p L$ for some $L' \in$ *NP–complete*, then L is *NP–hard* [4]. Moreover, if $L \in NP$, then $L \in$ *NP–complete* [4]. A principal *NP–complete* problem is *HAM–CYCLE* [4].

A simple graph is an undirected graph without multiple edges or loops [4]. An instance of the language *HAM–CYCLE* is a simple graph $G = (V, E)$ where V is the set of vertices and E is the set of edges, each edge being an unordered pair of vertices [4]. We say $(u, v) \in E$ is an edge in a simple graph $G = (V, E)$ where u and v are vertices. For a simple graph $G = (V, E)$, a simple cycle in G is a sequence of distinct vertices $\langle v_0, v_1, v_2, ..., v_k\rangle$ such that $(v_k, v_0) \in E$ and $(v_{i-1}, v_i) \in E$ for $i = 1, 2, ..., k$ [4]. A Hamiltonian cycle is a simple cycle of the simple graph which contains all the vertices of the graph. A simple graph that contains a hamiltonian cycle is said to be hamiltonian; otherwise, it is nonhamiltonian [4]. The problem *HAM–CYCLE* asks whether a simple graph is hamiltonian [4].

3 Results

▶ **Definition 1.** *A dense language on m is a formal language (a set of **binary** strings) where there exists a positive integer n_0 such that the counting of the number of strings of length $n \geq n_0$ in the language is greater than or equal to 2^{n-m} where m is a real number and $0 < m \leq 1$. The complexity class of all dense languages on m is called $DENSE(m)$.*

▶ **Definition 2.** *A formal language (a set of **binary** strings) is in $DENSE(0)$ if for every possible value of $0 < m \leq 1$, then the language is always in $DENSE(m)$.*

In this work, we are going to represent the simple graphs with an adjacency-matrix [4]. For the adjacency-matrix representation of a simple graph $G = (V, E)$, we assume that the vertices are numbered $1, 2, \ldots, |V|$ in some arbitrary manner. The adjacency-matrix representation of a simple graph G consists of a $|V| \times |V|$ matrix $A = (a_{i,j})$ such that $a_{i,j} = 1$ when $(i, j) \in E$ and $a_{i,j} = 0$ otherwise [4]. In this way, every simple graph of k vertices could be represented by a binary string of k^2 bits.

Observe the symmetry along the main diagonal of the adjacency matrix in this kind of graph that is called simple. We define the transpose of a matrix $A = (a_{i,j})$ to be the matrix $A^T = (a^T_{i,j})$ given by $a^T_{i,j} = a_{j,i}$. Hence the adjacency matrix A of a simple graph is its own transpose $A = A^T$.

▶ **Definition 3.** *The language NON–SIMPLE contains all the graph that are represented by an adjacency-matrix A such that $A \neq A^T$ or there is some $a_{i,j} = 1$ where $i = j$.*

▶ **Lemma 4.** *NON–SIMPLE $\in P$.*

Proof. Given a binary string x, we can check whether x is an adjacency-matrix which is not equal to its own transpose in time $O(|x|^2)$ just iterating each bit $a_{i,j}$ in x and checking whether $a_{i,j} \neq a_{j,i}$ or $a_{i,j} = 1$ when $i = j$ where $|\ldots|$ represents the bit-length function [4]. ◀

▶ **Definition 5.** *The language HAM–CYCLE' contains all the binary strings z such that $z = xy$, the bit-length of x is equal to $(\lfloor \sqrt{|z|} \rfloor)^2$ and $x \in$ HAM–CYCLE or $x \in$ NON–SIMPLE where y could be the empty string when $|\ldots|$ and $\lfloor \ldots \rfloor$ represent the bit-length function and the floor function respectively.*

▶ **Lemma 6.** *HAM–CYCLE' $\in$ NP–complete.*

Proof. Given a binary string z such that $z = xy$ and the bit-length of x is equal to $(\lfloor \sqrt{|z|} \rfloor)^2$, we can decide in polynomial time whether $x \notin$ *NON–SIMPLE* just verifying when $x = x^T$ and $a_{i,i} = 0$ for all vertex i. In this way, we can reduce in polynomial time a simple graph $G = (V, E)$ of k vertices encoded as the binary string x such that when x has k^2 bits and $x \notin$ *NON–SIMPLE* then

$x \in$ HAM–CYCLE if and only if $xy \in$ HAM–CYCLE'

where y could be the empty string. In this way, we can reduce in polynomial time each element of *HAM–CYCLE* to some element of *HAM–CYCLE'*. Therefore, *HAM–CYCLE'* is in *NP–hard*. Moreover, we can check in polynomial time over a binary string z such that $z = xy$ and the bit-length of x is equal to $(\lfloor \sqrt{|z|} \rfloor)^2$ whether $x \in$ *HAM–CYCLE* or $x \in$ *NON–SIMPLE* since *HAM–CYCLE* $\in NP$ and *NON–SIMPLE* $\in NP$ because of $P \subseteq NP$ [8]. Consequently, *HAM–CYCLE'* is in *NP*. Hence, *HAM–CYCLE'* $\in$ *NP–complete*. ◀

▶ **Lemma 7.** *HAM–CYCLE'* $\in DENSE(1)$. *This would mean the existence of a sufficiently large positive integer* n'_0 *such that all the binary strings of length* $n \geq n'_0$ *which belong to HAM–CYCLE' are more than or equal to* 2^{n-1} *elements.*

Proof. *OEIS A000088* gives some number of graphs on n unlabeled points [10]. For 8 points there are 12346 so just over half the graphs on 8 points are Hamiltonian [10]. For 12 points, there are 152522187830 Hamiltonian graphs out of 165091172592 which would claim that over 92% of the 12 point graphs are Hamiltonian [10]. For $n = 2$ there are two graphs, neither of which is Hamiltonian [10]. For $n < 8$ over half the graphs are not Hamiltonian [10]. It does not seem surprising that once n gets large most graphs are Hamiltonian [10].

Choosing a graph on n vertices at random is the same as including each edge in the graph with probability $\frac{1}{2}$, independently of the other edges [2]. You get a more general model of random graphs if you choose each edge with probability p [2]. This model is known as $G_{n,p}$ [2]. It turns out that for any constant $p > 0$, the probability that $G_{n,p}$ contains a Hamiltonian cycle tends to 1 when n tends to infinity [2]. In fact, this is true whenever $p > \frac{c \times \log n}{n}$ for some constant c. In particular this is true for $p = \frac{1}{2}$, which is our case [2].

For all the binary strings z such that $z = xy$ and the bit-length of x is equal to $(\lfloor \sqrt{|z|} \rfloor)^2$, the amount of elements of size $|z|$ in *HAM–CYCLE'* is equal to the number of binary strings $x \in$ *HAM–CYCLE* or $x \in$ *NON–SIMPLE* of size $(\lfloor \sqrt{|z|} \rfloor)^2$ multiplied by $2^{|z| - (\lfloor \sqrt{|z|} \rfloor)^2}$. Since the number of Hamiltonian graphs increases as much as we go further on n, it does not seem surprising either that once n gets large most binary strings belong to *HAM–CYCLE'*. Moreover, the amount of binary strings which have some bit-length k^2 and belongs to *NON–SIMPLE* is considerably superior to the amount of strings with the same bit-length which are valid simple graphs. Actually, we can affirm for a sufficiently large positive integer n'_0, all the binary strings of length $n \geq n'_0$ which belong to *HAM–CYCLE'* are indeed more than or equal to 2^{n-1} elements. In this way, we show that *HAM–CYCLE'* $\in DENSE(1)$. ◀

▶ **Definition 8.** *We will define a sequence of languages HAM–CYCLE'*$_k$ *for every possible integer* $1 \leq k$. *We state HAM–CYCLE'*$_1$ *as the language HAM–CYCLE'. Recursively, from a language HAM–CYCLE'*$_k$*, we define HAM–CYCLE'*$_{k+1}$ *as follows: A binary string* xy *complies with* $xy \in$ *HAM–CYCLE'*$_{k+1}$ *if and only if* x *and* y *are binary strings,* $x \in$ *HAM–CYCLE'*$_k$ *or* $y \in$ *HAM–CYCLE'*$_k$ *such that* $|x| = \lfloor \frac{|xy|}{2} \rfloor$ *where* $|\ldots|$ *represents the bit-length function and* $\lfloor \ldots \rfloor$ *is the floor function.*

▶ **Lemma 9.** *For every integer* $1 \leq k$*, HAM–CYCLE'*$_k$ $\in NP$.

Proof. This is true for $k = 1$ as we see in Lemma 6. Every string xy which belongs to *HAM–CYCLE'*$_2$ complies with $x \in$ *HAM–CYCLE'*$_1$ or $y \in$ *HAM–CYCLE'*$_1$ such that $|x| = \lfloor \frac{|xy|}{2} \rfloor$. Moreover, every string xy which belongs to the language *HAM–CYCLE'*$_3$ complies with $x \in$ *HAM–CYCLE'*$_2$ or $y \in$ *HAM–CYCLE'*$_2$ such that $|x| = \lfloor \frac{|xy|}{2} \rfloor$. Furthermore, we can extend this property for every positive integer $k > 3$ in *HAM–CYCLE'*$_k$. Indeed, *HAM–CYCLE'*$_k$ is in NP for every integer $1 \leq k$, since the verification of whether the two substrings are indeed elements of *HAM–CYCLE'*$_{k-1}$ can be done in polynomial time with the appropriated certificates using the induction on k. ◀

▶ **Theorem 10.** *For every integer* $1 \leq k$*, HAM–CYCLE'*$_k$ $\in$ *NP–complete.*

Proof. This is true for $k = 1$ by the Lemma 6. Let's assume it is valid for some positive integer $1 \leq k'$. Let's prove this for $k' + 1$. We already know the adjacency-matrix of n^2 zeros represents a simple graph of n vertices which does not contain any edge. This kind of a simple graph does not belong to *HAM–CYCLE'*$_1$. As a consequence, this string will not belong to any *HAM–CYCLE'*$_{k'}$, because its substrings of a quadratic length are also adjacency-matrix of only zeros. Suppose, we have an instance y of *HAM–CYCLE'*$_{k'}$. We can reduce y in *HAM–CYCLE'*$_{k'}$ to zy in *HAM–CYCLE'*$_{k'+1}$ such that

$$y \in \textit{HAM–CYCLE'}_{k'} \textit{ if and only if } zy \in \textit{HAM–CYCLE'}_{k'+1}$$

where the binary string z is exactly a sequence of $\lfloor \frac{|zy|}{2} \rfloor$ zeros. We can do this since we already know $z \notin$ *HAM–CYCLE'*$_{k'}$. Certainly, if the membership $zy \in$ *HAM–CYCLE'*$_{k'+1}$ is true, $z \notin$ *HAM–CYCLE'*$_{k'}$ and $|z| = \lfloor \frac{|zy|}{2} \rfloor$, then $y \in$ *HAM–CYCLE'*$_{k'}$ also holds according to the Definition 8. Since this reduction remains in polynomial time for every positive integer $1 \leq k'$, then we show that *HAM–CYCLE'*$_{k'+1}$ is in *NP–hard*. Moreover, *HAM–CYCLE'*$_{k'+1}$ is also in *NP–complete*, because of the Lemma 9. ◀

▶ **Theorem 11.** *For every integer $1 \leq k$, if the language HAM–CYCLE'$_k$ is in $DENSE(k')$ for every instance of bit-length $n' \geq n_0$, then HAM–CYCLE'$_{k+1}$ is in $DENSE(\frac{k'}{2})$ for every instance of bit-length $n' \geq 2 \times n_0$.*

Proof. If the language *HAM–CYCLE'*$_k$ is in $DENSE(k')$ for every instance of bit-length $n' \geq n_0$, then for every integer $n \geq n_0$ the amount of elements of size $n+i$ in *HAM–CYCLE'*$_{k+1}$ (where $i \geq n_0$ and $i = \lfloor \frac{n+i}{2} \rfloor$) is greater than or equal to

$$2^{i-k'} \times 2^n + 2^{n-k'} \times (2^i - 2^{i-k'}).$$

This is because there must be more than or equal to $2^{i-k'}$ elements of size i in *HAM–CYCLE'*$_k$ which are prefixes of the binary strings of size $n+i$ in the language *HAM–CYCLE'*$_{k+1}$. We multiply that amount by 2^n since this is the number of different combinations of suffixes with length n in the binary strings of size $n+i$. Moreover, there must be more than or equal to $2^{n-k'}$ elements of size n in *HAM–CYCLE'*$_k$ which are suffixes of the binary strings of size $n+i$ in *HAM–CYCLE'*$_{k+1}$. We multiply that amount by $(2^i - 2^{i-k'})$ since this is the number of different combinations of prefixes with length i in the binary strings of size $n+i$ just avoiding to count the previous prefixes twice. If we join both properties, we obtain the sum described by the formula above.

Indeed, this formula can be simplified to

$$2^{n+i-k'} + 2^{n+i-k'} \times (2^0 - 2^{-k'})$$

and extracting a common factor we obtain

$$2^{n+i-k'} \times (1 + (1 - 2^{-k'}))$$

which is equal to

$$2^{n+i-k'} \times (2 - \frac{1}{2^{k'}}).$$

Nevertheless, for every real number $0 < k' \leq 1$ we have that

$$(2 - \frac{1}{2^{k'}}) \geq 2^{\frac{k'}{2}}.$$

Certainly, if we multiply both member of the inequality by $2^{k'}$, we obtain

$$(2^{k'+1} - 1) \geq 2^{k'+\frac{k'}{2}}$$

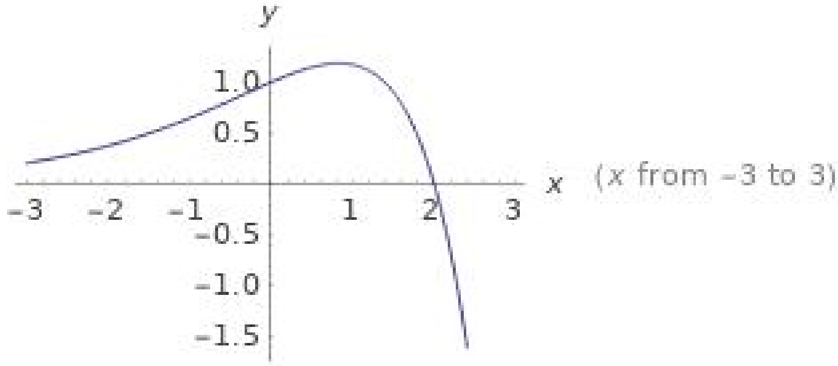

Figure 1 Plot the function f(x) on the interval [-3, 3]

which is equivalent to

$$2^{k'} \times (2 - 2^{\frac{k'}{2}}) \geq 1$$

that it is true for every real number $0 < k' \leq 1$. We can check in the Figure 1 that the function $f(x) = 2^x \times (2 - 2^{\frac{x}{2}})$ is greater than or equal to 1 over the interval $[0, 1]$. Thus

$$2^{n+i-k'} \times (2 - \frac{1}{2^{k'}}) \geq 2^{n+i-k'} \times 2^{\frac{k'}{2}}$$

where

$$2^{n+i-k'} \times 2^{\frac{k'}{2}} = 2^{n+i-(k'-\frac{k'}{2})} = 2^{n+i-\frac{k'}{2}}.$$

Since there are more than or equal to $2^{n'-(\frac{k'}{2})}$ elements of the language *HAM–CYCLE'*$_{k+1}$ with length $n' \geq 2 \times n_0$ therefore, we show that *HAM–CYCLE'*$_{k+1}$ is in $DENSE(\frac{k'}{2})$ for every instance of bit-length $n' \geq 2 \times n_0$. ◀

▶ **Lemma 12.** *HAM–CYCLE'*$_k$ $\in DENSE(\frac{1}{2^{k-1}})$ *for every instance of bit-length* $n \geq 2^{k-1} \times n'_0$*, where the constant* n'_0 *is the positive integer used in the Definition 1 and Lemma 7 for HAM–CYCLE'.*

Proof. According to the Lemma 7, *HAM–CYCLE'*$_1$ is in $DENSE(1)$ for every instance of bit-length $n \geq 2^0 \times n'_0 = n'_0$. Consequently, due to Theorem 11, *HAM–CYCLE'*$_2$ is in $DENSE(\frac{1}{2})$ for every instance of bit-length $n \geq 2^1 \times n'_0$. Moreover, *HAM–CYCLE'*$_3$ is in $DENSE(\frac{1}{4})$ for every instance of bit-length $n \geq 2^2 \times n'_0$ and so forth ... and thus, for every language *HAM–CYCLE'*$_k$, we have that *HAM–CYCLE'*$_k$ $\in DENSE(\frac{1}{2^{k-1}})$ for every instance of bit-length $n \geq 2^{k-1} \times n'_0$. ◀

▶ **Definition 13.** *We will define a language HAM–CYCLE'*$_\infty$ *as follows: A binary string* x *complies with* $x \in$ *HAM–CYCLE'*$_\infty$ *if and only if we obtain that* $x \in$ *HAM–CYCLE'*$_k$ *and* $2^{k-1} \times n'_0 \leq |x| < 2^k \times n'_0$ *where* $|\ldots|$ *represents the bit-length function and the constant* n'_0 *is the positive integer used in the Definition 1 and Lemma 7 for HAM–CYCLE'.*

▶ **Lemma 14.** *HAM–CYCLE'*$_\infty$ $\in NP$.

Proof. We can calculate the value of k from some binary string x that is approximately $\lceil \log_2(\frac{|x|}{n'_0}) \rceil$, where $\lceil \ldots \rceil$ is the ceiling function. In this way, we should know if

$x \in$ *HAM–CYCLE'*$_\infty$, then $x \in$ *HAM–CYCLE'*$_k$. However, for every positive integer k, we can check in polynomial time whether $x \in$ *HAM–CYCLE'*$_k$ just splitting the binary string x into the following substrings $x = x_1 x_2 x_3 \ldots x_{2^{k-1}}$ and verifying later whether $x_1 \in$ *HAM–CYCLE'*$_1$ or $x_2 \in$ *HAM–CYCLE'*$_1$ or $x_3 \in$ *HAM–CYCLE'*$_1$ and so forth ... until we finally check whether $x_{2^{k-1}} \in$ *HAM–CYCLE'*$_1$ where 2^{k-1} is polynomially bounded by the bit-length string $|x|$. Indeed, the language *HAM–CYCLE'*$_\infty$ is in NP, because the verification of whether the whole string or a polynomially amount of substrings are indeed elements of *HAM–CYCLE'*$_1$ can be done in polynomial time with the appropriated certificates. ◀

▶ **Theorem 15.** *HAM–CYCLE'*$_\infty \in$ *NP–complete.*

Proof. We already know the adjacency-matrix of n^2 zeros represents a simple graph of n vertices which does not contain any edge. This kind of a simple graph does not belong to *HAM–CYCLE'*$_1$. Suppose, we have an instance y of *HAM–CYCLE'*$_1$. We can reduce y in *HAM–CYCLE'*$_1$ to zy in *HAM–CYCLE'*$_\infty$ such that

$y \in$ *HAM–CYCLE'*$_1$ *if and only if* $zy \in$ *HAM–CYCLE'*$_\infty$

where z is a binary string of a sequence of zeros such that $2^{k-1} \times n'_0 \leq |zy| < 2^k \times n'_0$ and the membership in $zy \in$ *HAM–CYCLE'*$_k$ implies that $y \in$ *HAM–CYCLE'*$_1$, where the constant n'_0 is the positive integer used in the Definition 1 and Lemma 7 for *HAM–CYCLE'*. We claim that the bit-length of zy is polynomially bounded by $|y|$. Certainly, the bit-length of z is polynomially bounded by $2^{k-1} \times n'_0$ and $|y|$ since $k \approx \lceil \log_2(\frac{|zy|}{n'_0}) \rceil$, where $\lceil \ldots \rceil$ is the ceiling function. The previous expression would be equivalent to $2^k \approx \frac{|y| + 2^{k-1} \times n'_0}{n'_0}$ which means that $\frac{|y|}{2^k \times n'_0} \approx 1$. In this way, we show that *HAM–CYCLE'*$_\infty$ is in *NP–hard*. Moreover, we demonstrate that *HAM–CYCLE'*$_\infty$ is also in *NP–complete*, because of the Lemma 14. ◀

▶ **Lemma 16.** *HAM–CYCLE'*$_\infty \in DENSE(0)$.

Proof. When k tends to infinity, then $\frac{1}{2^{k-1}}$ tends to 0. In this way, we obtain that *HAM–CYCLE'*$_k \in DENSE(0)$ as a consequence of the Lemma 12. Actually, *HAM–CYCLE'*$_\infty$ contains the elements of the languages *HAM–CYCLE'*$_k$ into the interval of the binary strings between the bit-length $2^{k-1} \times n'_0 \leq n < 2^k \times n'_0$. Those elements will have a bit-length greater than $2^{k-1} \times n'_0$ and by the Lemma 12 the density in the interval would be $DENSE(\frac{1}{2^{k-1}})$. Therefore, the proof is done. ◀

4 Discussion

When a language is sparse, then its complement is in $DENSE(0)$ [6]. Indeed, the sparse languages are called sparse because there are a total of 2^n strings of length n, and if a language only contains polynomially many of these, then the proportion of strings of length n that it contains rapidly goes to zero as n grows (which means its complement should be in $DENSE(0)$) [6]. In addition, according to Theorem 15, the complement of this language *HAM–CYCLE'*$_\infty$ must be in *coNP-complete*, because of the complements of the *NP-complete* problems are complete for $coNP$ [8]. In 1999, Jin-Yi Cai and D. Sivakumar, building on work by Ogihara, showed that if there exists a sparse *P-complete* problem, then $LOGSPACE = P$ [3]. We might extend the proof of this paper to show the same result on P. Certainly, we might only need to find some *P-complete* which belongs to $DENSE(1)$ because the *P-completeness* is closed under complement [8]. Indeed, the other steps of that possible proof might be similar to the arguments that we follow in this paper.

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